#### Signature-based Gröbner basis algorithms over Tate algebras

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• Question: in  $\mathbb{R}[X]$ , reduce  $f = X^2$  modulo g = 0.01X - 1

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- ► The usual way:

$$f = X^{2}$$

$$\begin{pmatrix} -100Xg \\ 100X \\ -10\ 000g \\ 10\ 000 \end{pmatrix}$$

- It terminates, but...
- $g \simeq 1$ , but  $f \mod g \not\simeq 0$

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- Question: in  $\mathbb{R}[X]$ , reduce  $f = X^2$  modulo g = 0.000001X 1LT(g)
- ► The usual way:

 $f = X^{2}$   $\begin{pmatrix} -1\,000\,000Xg \\ 1\,000\,000X \\ (-1\,000\,000\,000\,000g \\ 1\,000\,000\,000\,000 \end{pmatrix}$ 

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Another way?  

$$f = X^{2}$$

$$\begin{pmatrix} +X^{2}g \\ 0.01X^{3} \\ (+0.01X^{3}g \\ 0.0001X^{4} \\ (\cdots$$

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- The sequence of reductions tends to 0

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- ► This work: make sense of this process for convergent power series in Z<sub>p</sub>[[X]]

- DVR = principal local domain  $K^{\circ}$  with maximal ideal  $\langle \pi \rangle$ , residue field  $\mathbb{F} = K^{\circ} / \langle \pi \rangle$ 
  - $\mathbb{Z}_p$ р Х  $\mathbb{F}_p$  $\mathbb{C}[[X]]$
- Elements can be written  $a = \sum_{n=0}^{\infty} a_n \pi^n$ ,  $a_n \in \mathbb{F}$
- Valuation of  $a = \max n$  such that  $\pi^n$  divides a
- $\mathbb{Z}_p$  and  $\mathbb{C}[[X]]$  are complete for this topology



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No loss of precision possible:
 if a and b are small, a + b is small

$$a+b=a+b$$



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#### **Tate Series**

#### Definition

K{X}° = ring of series in X with coefficients in K° converging for all x ∈ K°
 = ring of power series whose general coefficients tend to 0

#### Motivation

Introduced by Tate in 1971 for rigid geometry

(p-adic equivalent of the bridge between algebraic and analytic geometry over  $\mathbb{C}$ )

 $\mathbf{X} = X_1, \ldots, X_n$ 

#### Examples

Polynomials (finite sums are convergent)

• 
$$\sum_{i,j=0}^{\infty} \pi^{i+j} X^i Y^j = 1 + \pi X + \pi Y + \pi^2 X^2 + \pi^2 XY + \pi^2 Y^2 + \cdots$$
  
• Not a Tate series: 
$$\sum_{i=0}^{\infty} X^i = 1 + 1X + 1X^2 + 1X^3 + \cdots$$

# Term ordering for Tate algebras

- Starting from a usual monomial ordering  $1 <_m \mathbf{X}^{\mathbf{i}_1} <_m \mathbf{X}^{\mathbf{i}_2} <_m \dots$
- ▶ We define a term ordering putting more weight on large coefficients



 $\mathbf{X}^{\mathbf{i}} = X_1^{i_1} \cdots X_n^{i_n}$ 

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- It has infinite descending chains, but they converge to zero
- Tate series always have a leading term

$$f = a_2 XY + a_1 X + a_0 \cdot 1 + a_3 X^2 Y^2 + \dots$$

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lsomorphism 
$$K\{\mathbf{X}\}^{\circ}/\langle \pi \rangle \simeq \mathbb{F}[\mathbf{X}]$$
  
 $f \mapsto \overline{f}$ 

compatible with the term order

$$f = \overbrace{a_2 XY + a_1 X}^{\mathsf{LT}(f)} + \overbrace{a_0}^{\circ} \cdot 1 + \overbrace{a_3}^{\circ} X^2 Y^2 + \dots$$
  
$$\overline{f} = \overline{a_2 XY + \overline{a_1} X}$$

 $\mathbf{X}^{\mathbf{i}} = X_1^{i_1} \cdots X_n^{i_n}$ 

## Gröbner bases

Standard definition once the term order is defined:

*G* is a Gröbner basis of  $I \iff$  for all  $f \in I$ , there is  $g \in G$  s.t. LT(g) divides LT(f)

- Standard equivalent characterizations:
  - 1. G is a Gröbner basis of I
  - 2. for all  $f \in I$ , f is reducible modulo G
  - 3. for all  $f \in I$ , f reduces to zero modulo G

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 $\exists$  sequence of reductions converging to 0

$$\pi f \in I \implies f \in I$$

4.  $\overline{G}$  is a Gröbner basis of  $\overline{I}$  in the sense of  $\mathbb{F}[\mathbf{X}]$ 

### Gröbner bases

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#### If *I* is saturated:

- 4.  $\overline{G}$  is a Gröbner basis of  $\overline{I}$  in the sense of  $\mathbb{F}[\mathbf{X}]$
- Every Tate ideal has a finite Gröbner basis
- ▶ It can be computed using the usual algorithms (reduction, Buchberger, F₄)
- ► In practice, the algorithms run with finite precision and without loss of precision

 $\exists$  sequence of reductions converging to 0

$$\pi f \in I \implies f \in I$$

No division by  $\pi$ 

### What about Tate series over a field?

- ► CDVF = fraction field K of a CDVR  $K^{\circ}$   $\mathbb{Q}_{p}$   $\mathbb{Z}_{p}$  $\mathbb{C}((X))$   $\mathbb{C}[[X]]$
- Elements can be written  $a = \sum_{n=-r}^{\infty} a_n \pi^n$ ,  $a_n \in \mathbb{F}$
- The valuation can be negative but not infinite
- Same metric, same topology as  $K^{\circ}$



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- Elements can be written  $a = \sum_{n=-r}^{\infty} a_n \pi^n$ ,  $a_n \in \mathbb{F}$
- The valuation can be negative but not infinite
- Same metric, same topology as  $K^{\circ}$
- Tate series can be defined as in the integer case
- Same order, same definition of Gröbner bases
- Main difference: πX now divides X
- Another surprising equivalence
  - 1. G is a normalized GB of I
  - 2.  $G \subset K{\mathbf{X}}^\circ$  is a GB of  $I \cap K{\mathbf{X}}^\circ$



$$\forall g \in G, LC(g) = 1 \text{ (in part., } G \subset K\{\mathbf{X}\}^{\circ})$$

 In practice, we emulate computations in K{X}° in order to avoid losses of precision (and the ideal is saturated)

Problem: useless and redundant computations, infinite reductions to 0

Example with a S-polynomial

$$p = p_1 f_1 + p_2 f_2 + \dots + p_k f_k + \dots + p_m f_m \qquad q = q_1 f_1 + q_2 f_2 + \dots + q_l f_l + \dots + q_m f_m$$

 $\text{S-Pol}(p,q) = \mu p - \nu q$ 

Problem: useless and redundant computations, infinite reductions to 0

 1<sup>st</sup> idea: keep track of the representation of the ideal elements [Möller, Mora, Traverso 1992]

$$p = p_1f_1 + p_2f_2 + \dots + p_kf_k + \dots + p_mf_m \qquad q = q_1f_1 + q_2f_2 + \dots + q_lf_l + \dots + q_mf_m$$
$$\mathbf{p} = p_1\mathbf{e}_1 + p_2\mathbf{e}_2 + \dots + p_k\mathbf{e}_k + \dots + p_m\mathbf{e}_m \qquad \mathbf{q} = q_1\mathbf{e}_1 + q_2\mathbf{e}_2 + \dots + q_l\mathbf{e}_l + \dots + q_m\mathbf{e}_m$$

S-Pol
$$(p, q) = \mu p - \nu q$$
  
S-Pol $(\mathbf{p}, \mathbf{q}) = \mu (p_1 \mathbf{e}_1 + \dots + p_k \mathbf{e}_k + \dots + p_m \mathbf{e}_m) - \nu (q_1 \mathbf{e}_1 + \dots + q_l \mathbf{e}_l + \dots + q_m \mathbf{e}_m)$ 

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- 1<sup>st</sup> idea: keep track of the representation of the ideal elements [Möller, Mora, Traverso 1992]
- 2<sup>nd</sup> idea: the largest term of the representation is enough [Faugère 2002; Gao, Volny, Wang 2010; Arri, Perry 2011... Eder, Faugère 2017]

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$$= \text{smaller terms} + \text{LT}(p_k) \mathbf{e}_k$$

$$= \text{smaller terms} + \text{LT}(q_l) \mathbf{e}_l$$

S-Pol(p, q) = 
$$\mu p - \nu q$$
  
S-Pol(p, q) =  $\mu (p_1 \mathbf{e}_1 + \dots + p_k \mathbf{e}_k + \dots + 0 \mathbf{e}_m) - \nu (q_1 \mathbf{e}_1 + \dots + q_l \mathbf{e}_l + \dots + 0 \mathbf{e}_m)$   
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= smaller terms +  $\mu LT(p_k) \mathbf{e}_k$  if  $\mu LT(p_k) \mathbf{e}_k \ge \nu LT(q_l) \mathbf{e}_l$  Regular S-polynomia

# Signatures for Tate algebra

#### Main properties of signatures:

- Ordered (in a way compatible with monomials)
- ► Example: Position over Term:  $\mu \mathbf{e}_i < \nu \mathbf{e}_j \iff i < j$  or i = j and  $\mu < \nu$
- Never decreasing in the course of the algorithms

#### Difficulties with Tate series:

- Need to order them with their coefficients
- The order is mixed:  $1 > \pi$

#### **Results**:

- Proof of correctness and termination for two orders:
  - Position over Term
  - > Valuation over Position over Term: analogue of the F5 order for the valuation
- No need to multiply signatures by  $\pi$

# Conclusion and perspectives

#### Main results

- Definitions of Gröbner bases for Tate series
- Algorithms for computing and using those Gröbner bases
- Data structure and algorithms implemented in Sage (version 8.5, 22/12/2018)
- Two signature-based algorithms with significant performance improvements

#### Perspectives

- Reduction of Tate series is very different from reduction of polynomials
- Design algorithms to perform those reductions more efficiently
- ► Goal: being able to take advantage of e.g. delaying reductions using signatures

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# Thank you for your attention!

#### More information and references:

- Xavier Caruso, Tristan Vaccon and Thibaut Verron (2019). 'Gröbner bases over Tate algebras'. In: ISSAC'19, arXiv:1901.09574. arXiv: 1901.09574 [math.AG]
- Xavier Caruso, Tristan Vaccon and Thibaut Verron (Feb. 2020). 'Signature-based algorithms for Gröbner bases over Tate algebras'. In: URL: https://hal.archives-ouvertes.fr/hal-02473665

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 $\overset{\bullet}{f} - \overset{\bullet}{q_1} \overset{\bullet}{\overline{g_1}} - \overset{\bullet}{q_2} \overset{\bullet}{\overline{g_2}} - \dots - \overset{\bullet}{q_r} \overset{\bullet}{\overline{g_r}} = 0$ 

4. So we have a sequence of reductions



1. Start with  $f \in I$ , we can assume that f has valuation 0

2. Separate  $f = \overline{f} + f - \overline{f}$ 

3.  $\overline{f} \in \overline{I}$  so we have a sequence of reductions

 $\mathbf{\hat{f}} - \mathbf{q}_1 \mathbf{\hat{g}}_1 - \mathbf{q}_2 \mathbf{\hat{g}}_2 - \dots - \mathbf{q}_r \mathbf{\hat{g}}_r = 0$ 

4. So we have a sequence of reductions

$$f - \sum_{i=1}^{r} \mathbf{q}_{i} g_{i} = f - \sum_{i=1}^{r} \mathbf{q}_{i} \overline{g}_{i} + \sum_{i=1}^{r} \mathbf{q}_{i} \left( \frac{\mathbf{e}}{g_{i}} - g_{i} \right)$$

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- 1. Start with  $f \in I$ , we can assume that f has valuation 0
- 2. Separate  $f = \overline{f} + f \overline{f}$

3.  $\overline{f} \in \overline{I}$  so we have a sequence of reductions

 $\mathbf{\hat{f}} - \mathbf{q}_1 \mathbf{\hat{g}}_1 - \mathbf{q}_2 \mathbf{\hat{g}}_2 - \dots - \mathbf{q}_r \mathbf{\hat{g}}_r = 0$ 

4. So we have a sequence of reductions

$$f - \sum_{i=1}^{r} \mathbf{q}_{i} g_{i} = f - \sum_{i=1}^{r} \mathbf{q}_{i} \overline{g_{i}} + \sum_{i=1}^{r} \mathbf{q}_{i} \left( \mathbf{e}_{i} - g_{i} \right)$$
$$= f - \overline{f} + \sum_{i=1}^{r} \mathbf{q}_{i} \left( \overline{g_{i}} - g_{i} \right)$$

I is saturated

1. Start with  $f \in I$ , we can assume that f has valuation 0

2. Separate  $f = \overline{f} + f - \overline{f}$ 

3.  $\overline{f} \in \overline{I}$  so we have a sequence of reductions

 $\mathbf{\hat{f}} - \mathbf{q}_1 \mathbf{\hat{g}}_1 - \mathbf{q}_2 \mathbf{\hat{g}}_2 - \dots - \mathbf{q}_r \mathbf{\hat{g}}_r = 0$ 

4. So we have a sequence of reductions

$$f - \sum_{i=1}^{r} \mathbf{q}_{i} \mathbf{g}_{i} = f - \sum_{i=1}^{r} \mathbf{q}_{i} \mathbf{\overline{g}}_{i} + \sum_{i=1}^{r} \mathbf{q}_{i} \left( \mathbf{\overline{g}}_{i} - \mathbf{g}_{i} \right)$$
$$= f - \overline{f} + \sum_{i=1}^{r} \mathbf{q}_{i} \left( \mathbf{\overline{g}}_{i} - \mathbf{g}_{i} \right) = \mathbf{\overline{m}} = \pi \cdot f_{1}$$

I is saturated

1. Start with  $f \in I$ , we can assume that f has valuation 0 2. Separate  $f = f + f - \overline{f}$ 3.  $\overline{f} \in \overline{I}$  so we have a sequence of reductions  $\overline{f} - q_1 \overline{g_1} - q_2 \overline{g_2} - \cdots - q_r \overline{g_r} = 0$ 4. So we have a sequence of reductions  $f - \sum_{i=1}^{r} \mathbf{q}_{i} \mathbf{g}_{i} = f - \sum_{i=1}^{r} \mathbf{q}_{i} \mathbf{\overline{g}}_{i} + \sum_{i=1}^{r} \mathbf{q}_{i} \left( \mathbf{\overline{g}}_{i} - \mathbf{g}_{i} \right)$  $= f - \overline{f} + \sum_{i=1}^{r} q_i \left(\overline{g_i} - g_i\right) = \stackrel{\circ}{\blacksquare} = \pi \cdot f_1$ 

*I* is saturated